

# Auto-Approximation of Graph Computing Zechao Shang, Jeffrey Xu Yu The Chinese University of Hong Kong

## Motivation

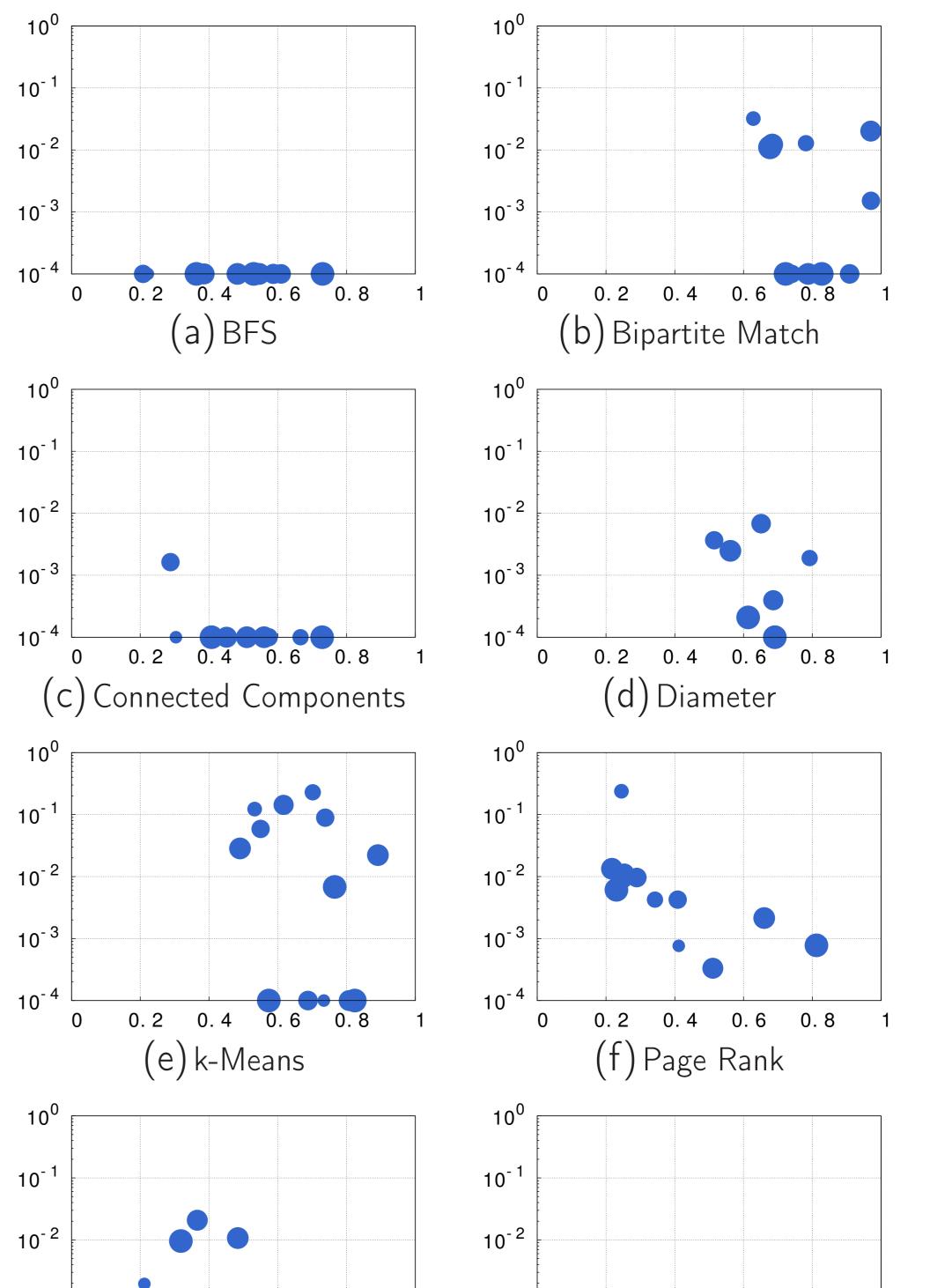
- Why Approximate?
  - ► **Big** data: GB, TB, PB, EB, and still increasing.
  - **Slow** algorithm: from  $O(n^2)$  to NP-Complete everywhere.
  - Approximate or not approximate?
    - Big data beats clever algorithm.
    - Approximated answers are usually sufficient for analytics.
    - ► Approximated good vs. exact bad. Which will you choose?
  - Answer delayed is answer denied.
- Why Auto?
  Developing "big data" solution is hard. That's why we are having this conference.
   Designing approximation algorithms is hard.
   Designing distributed approx. algorithms is hard + hard.

#### Does it work?

We conducted experiments on a wide spectrum of algorithms and datasets.

They demonstrates our approach works quite well on both convex and non-convex (!), continuous and discrete (!) cases.

x-axis is relative computation time compared to the original (un-approx.) algorithm, y-axis is relative error (in *log* scale), dots for datasets. Dot size for data size.



- Our Target
  - Minimum level of user involvement.
  - ► Approximation "for dummies".
  - ► Ultimate Goal: "One button" approximation.

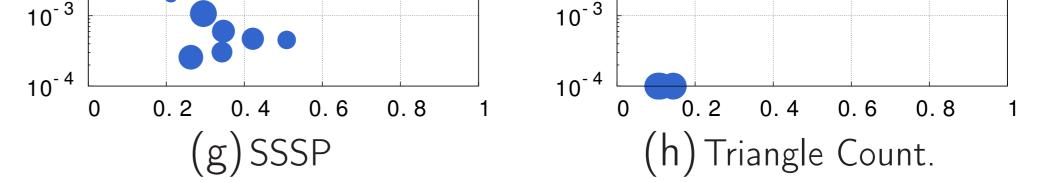
# Our approach

# Sample the **PROGRAM**, not the **GRAPH**!

- ► Why graph sampling is NOT working:
  - ► The vertices and edges are not independent at all.
  - ► The "connections" are exactly what we care most.
- ► How we do it?
  - Background: iterative, vertex-centric graph analytics.
    - ► We ignore some of the program instructions. That's it!
    - And carefully calibrate the answer to keep it unbiased.
      (See below.)
- Why it works?
  - Can we always do it?
    - Mostly. Even for complex queries, their vertex-centric instructions are usually simple and regular.
  - ► Is the error small?
    - Mostly. Under mild assumptions, we show the error will converge (to a small value).
    - ► We analyze it by analogizing it to the *social network information diffusion* and denote it as *error drifting*. Check paper for details.

#### **Toolbox for Program Sampling**

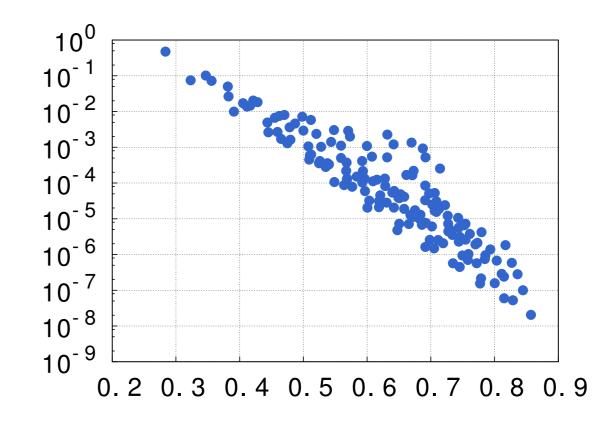
For example, if the user sums all incoming messages, we sample each five of them.



### **Opportunities and Challenges**

By manipulating parameters it generates a wide spread trade-off between accuracy and time. We are investigating how to make full use of it. (Cost model? Discussions are welcomed.)

x-axis is relative computation time compared to the original (un-approx.) algorithm, y-axis is relative error (in *log* scale), dots for combinations of parameters.



Works for both synchronize and asynchronize computing. We will try distributed environment later.

- This is not graph sampling!
- ► We find such opportunities using program analyzing.
- And compensate the final sum to keep it unbiased.
- Other alternatives:
  - Memorization: remember past answers.
  - ► Task skipping: take a nap this time. Work next call.
  - Interpolation: simple function replace complex one.
  - System function replacement: 1 + x for exp(x).
  - And still extending!

